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NESA Lab

Detecting Missed Security Operations Through Differential Checking of Object-based Similar Paths

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Background

Background

Reference count operation

Security check

Initialization

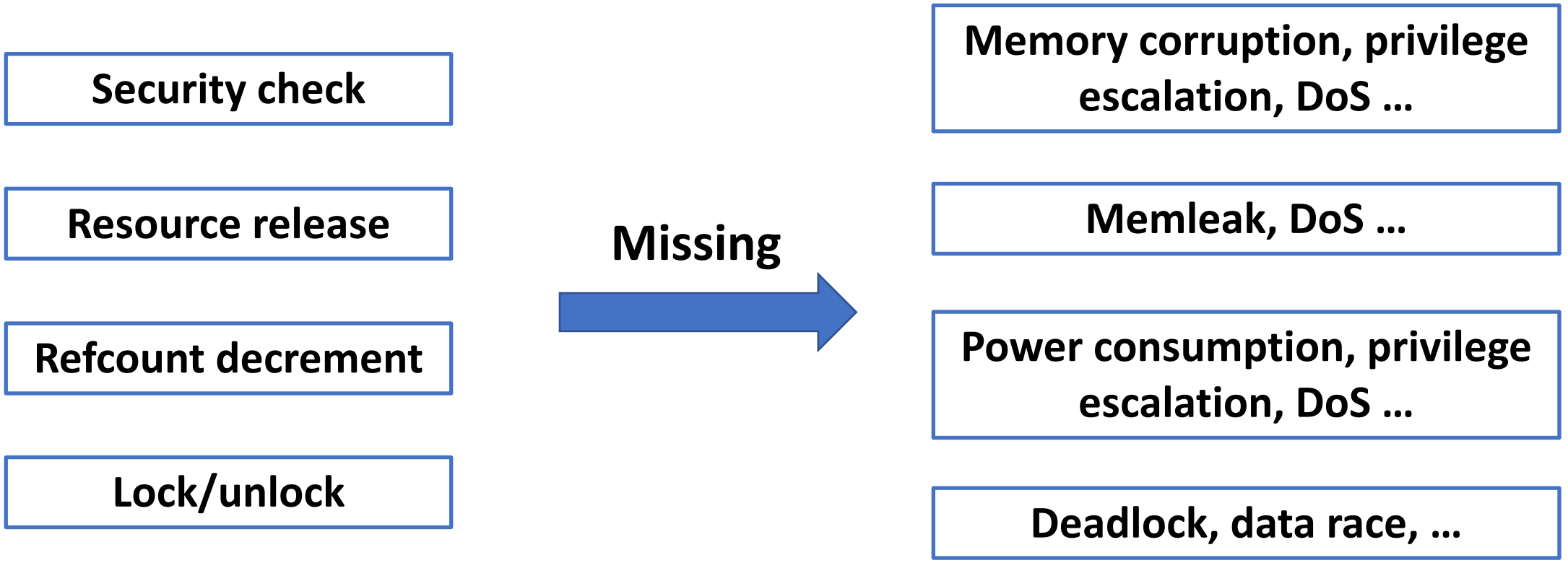
Security operations are widely used in large-scale programs

Lock

Resource release

Background

- Missing security operations could lead to many **security issues**



61% vulnerabilities in the NVD are caused by missing security operations!

Background

How to determine whether the missed security operations are indeed necessary ?

Cross-checking

- **High level idea**
  - Collect a substantial number of similar code pieces.
  - Check the behaviors of security operations across the similar code pieces.
  - The majority is correct.

- **Limitations**
  - Sufficient code pieces are required to enable cross-checking.
  - The granularity of code piece is hard to control.
  - The majority is not always correct.

Insight

- **A security operation usually focuses on one critical object.**

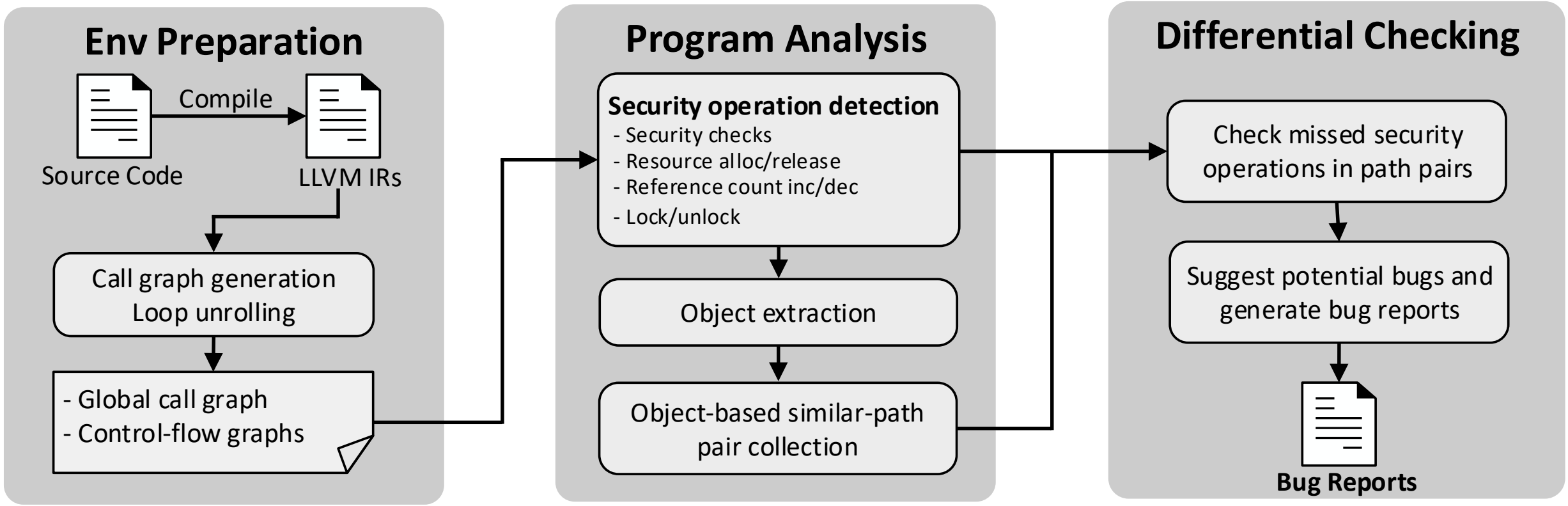
- **The similarity of code pieces should be based on the particular object.**
  - Object-based similar path pair.
  - It takes only 2 paths to enable inconsistency analysis and bug detection.
  - Fine-grained and robust.

System Design

Overview

IPPO (Inconsistent Path Pairs as a bug Oracle)

- Statically detect bugs caused by missed security operations.
- LLVM-based intra-procedural static analyzer.



Security Operation Detection

Security check

FILE: drivers/dma/dma-jz4780.c

```
…
854.    jzdma = devm_kzalloc(dev, struct_size(jzdma, chan,
855.                         soc_data->nb_channels), GFP_KERNEL);
856.    if (!jzdma)
857.        return -ENOMEM;
…
```

Lock/unlock

FILE: arch/x86/platform/uv/uv_irq.c

```
…
161.    mutex_lock(&uv_lock);
…
175.    mutex_unlock(&uv_lock);
…
```

Refcount inc/dec

FILE: drivers/net/ethernet/intel/e1000e/ethtool.c

```
…
161.    pm_runtime_get_sync(netdev->dev.parent);
…
175.    pm_runtime_put_sync(netdev->dev.parent);
…
```

Resource alloc/release

FILE: drivers/platform/x86/dell/dell-wmi-sysman/biosattr-interface.c

```
…
124.    buffer = kzalloc(buffer_size, GFP_KERNEL);
…
141.    kfree(buffer);
…
```

Extracting Objects

Security check

FILE: drivers/dma/dma-jz4780.c

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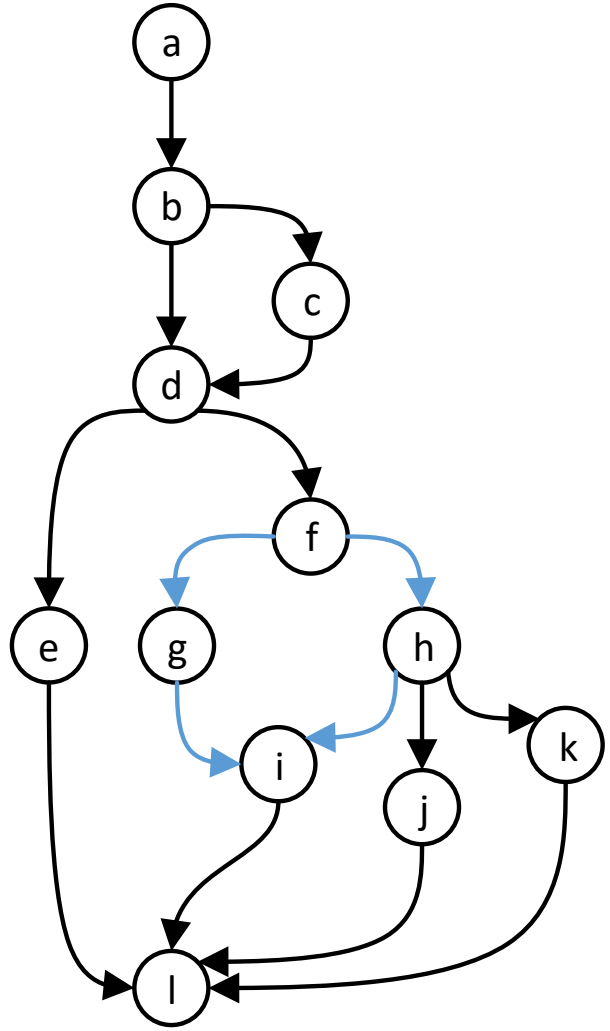
Object-based Similar Path Pair

➢ **Rules for constructing object-based similar path pair (OSPP)**

- **Rule 1**
  - The two paths start at the same block and end at the same block in CFG.

- **Challenge: path explosion in large functions**



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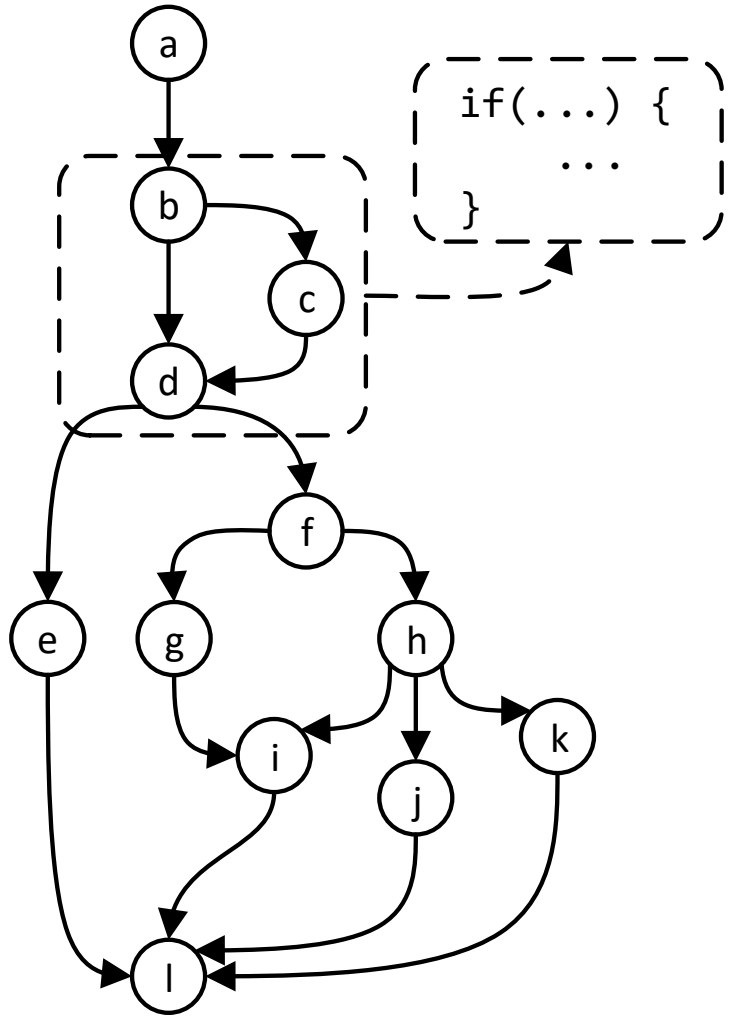
Root cause: The redundant common messages



Object-based Similar Path Pair

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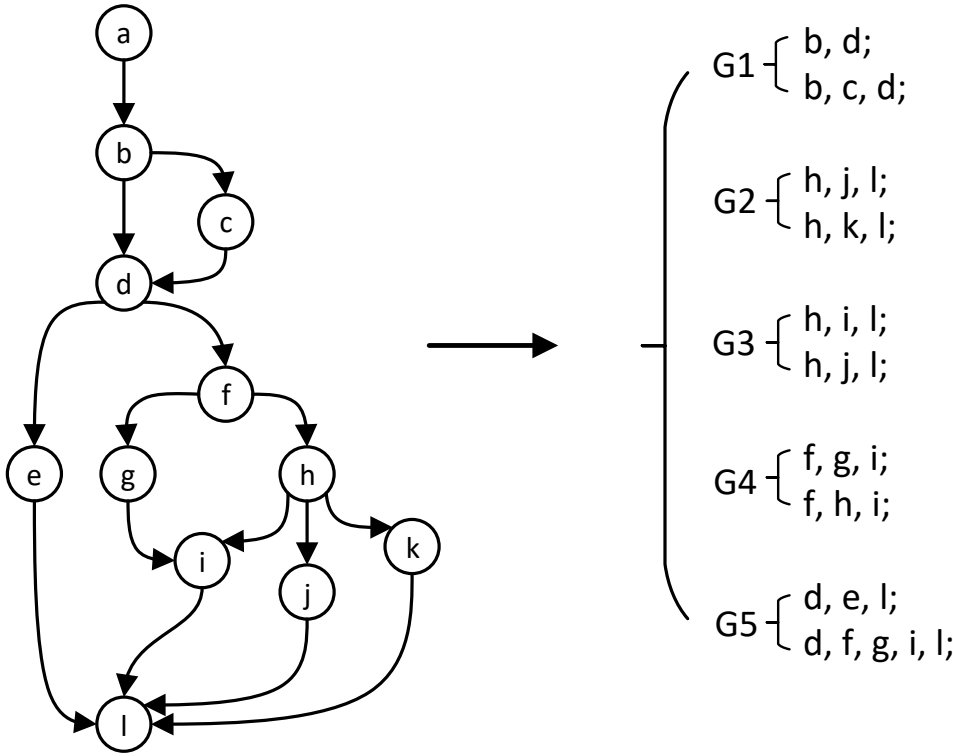
- **Rule 1**
  - The two paths start at the same block and end at the same block in CFG.

- **Challenge: path explosion in large functions**

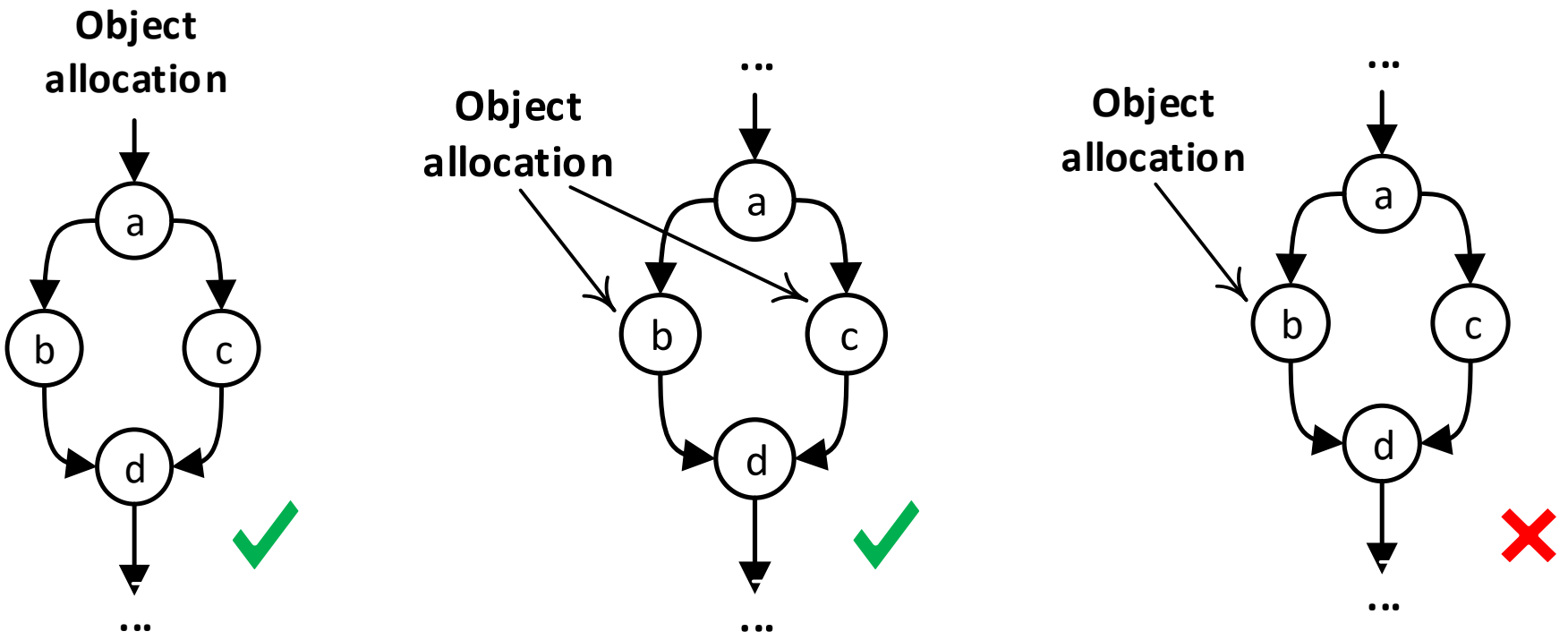
- **Our solution: reduced similar path (RSP)**
  - Only collect paths that share no common basic blocks besides the start block and the end block.



Object-based Similar Path Pair

➢ **Rules for constructing object-based similar path pair (OSPP)**

- **Rule 2**
  - The object has the same state in two paths.



Object-based Similar Path Pair

➢ **Rules for constructing object-based similar path pair (OSPP)**

- **Rule 3**
  - The object has the same ***security operation-influential operations*** against the object.

Table 1: SO-influential operations.

Security operation	SO-influential operation
Security check	Function calls, arithmetic and memory operations after the object (checked variable)
Resource alloc/release	Resource propagation
Refcount	Reference counter adjustment
Lock/unlock	Lock state adjustment

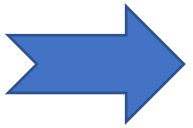
Object-based Similar Path Pair

➢ **Rules for constructing object-based similar path pair (OSPP)**

- **Rule 4**
  - The two paths have the same set of pre- and post-conditions against the object.

Pre-condition:

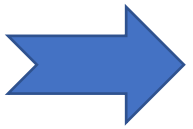
The branch condition of a path pair.



Must be object irrelevant

Post-condition:

The return values of a path pair.



A pair of normal paths

or a pair of error handing paths

Object-based Similar Path Pair

➢ **Rules for constructing object-based similar path pair (OSPP)**

- **Challenge:** how to efficiently collect path pairs that satisfy the post-condition of Rule 4?

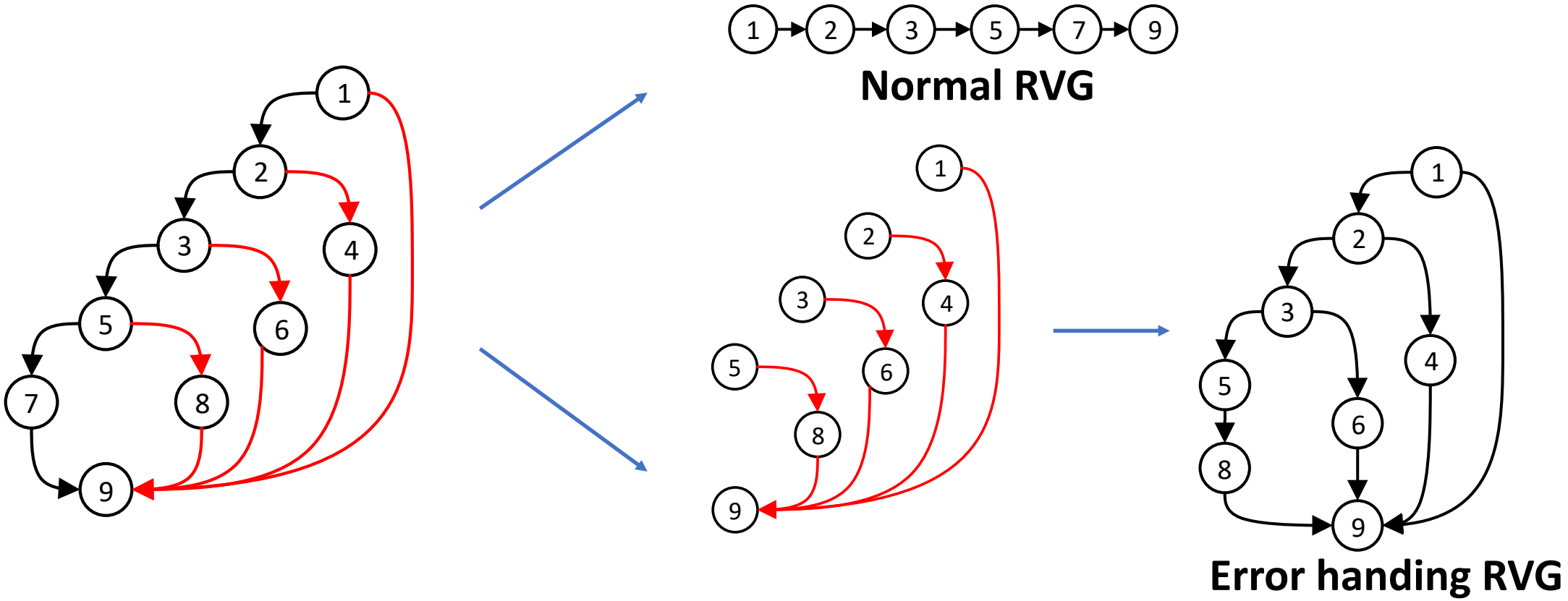
- **Our solution: graph partitioning**
  - Divide the CFG into 2 sub-CFGs:
  - Paths in each sub-CFG share the same return value

➢ Return value-based graphs (RVGs)

Object-based Similar Path Pair

- **Rules for constructing object-based similar path pair (OSPP)**

  - **Generating return value-based graphs**



Case Study

A Double-free Bug Found by IPPO

```
1  /* sound/pci/echoaudio/echoaudio.c */
2  static int snd_echo_resume(struct device *dev)
3  {
4      struct echoaudio *chip = dev_get_drvdata(dev);
5      struct comm_page *commpage, *commpage_bak;
6      ...
7      commpage = chip->comm_page;
8      commpage_bak = kmemdup(commpage, sizeof(*commpage), GFP_KERNEL);
9      if (commpage_bak == NULL)
10         return -ENOMEM;
11
12     err = init_hw(chip, chip->pci->device, chip->pci->subsystem_device);
13     if (err < 0) {
14         kfree(commpage_bak);
15         dev_err(dev, "resume init_hw err=%d\n", err);
16         snd_echo_free(chip);
17         return err;
18     }
19     ...
20     err = restore_dsp_rettings(chip);
21     chip->pipe_alloc_mask = pipe_alloc_mask;
22     if (err < 0) {
23         kfree(commpage_bak);
24         return err;
25     }
26     ...
27     kfree(commpage_bak);
28     ...
29     if(request_irq(...)) {
30         dev_err(chip->card->dev, "cannot grab irq\n");
31         snd_echo_free(chip);
32         return -EBUSY;
33     }
34     ...
35     return 0;
36 }
```

Resource allocation of object *chip*

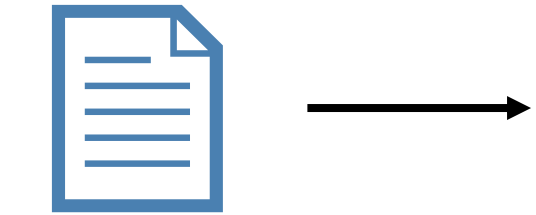
Resource release of object *chip*

Missing release against *chip* in the error handing path

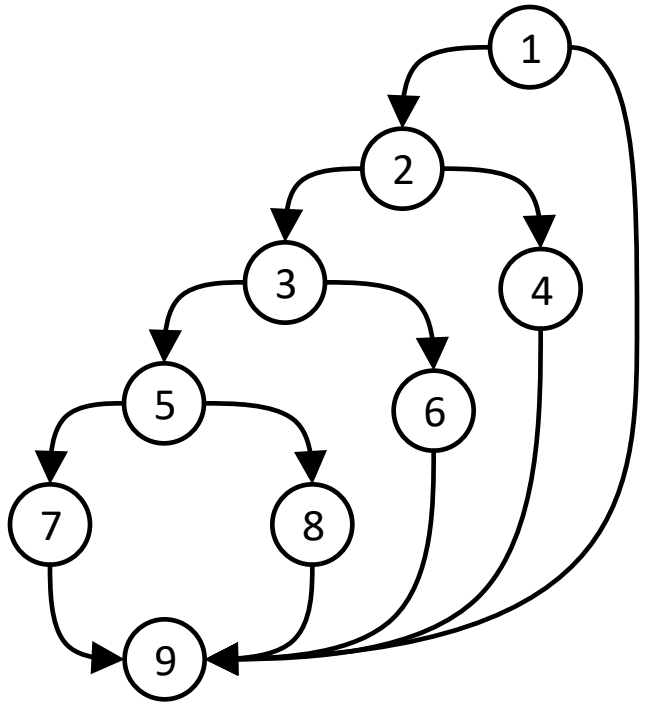
Resource release of object *chip*

Workflow of IPPO

- **Security operation detection & error edges identification**



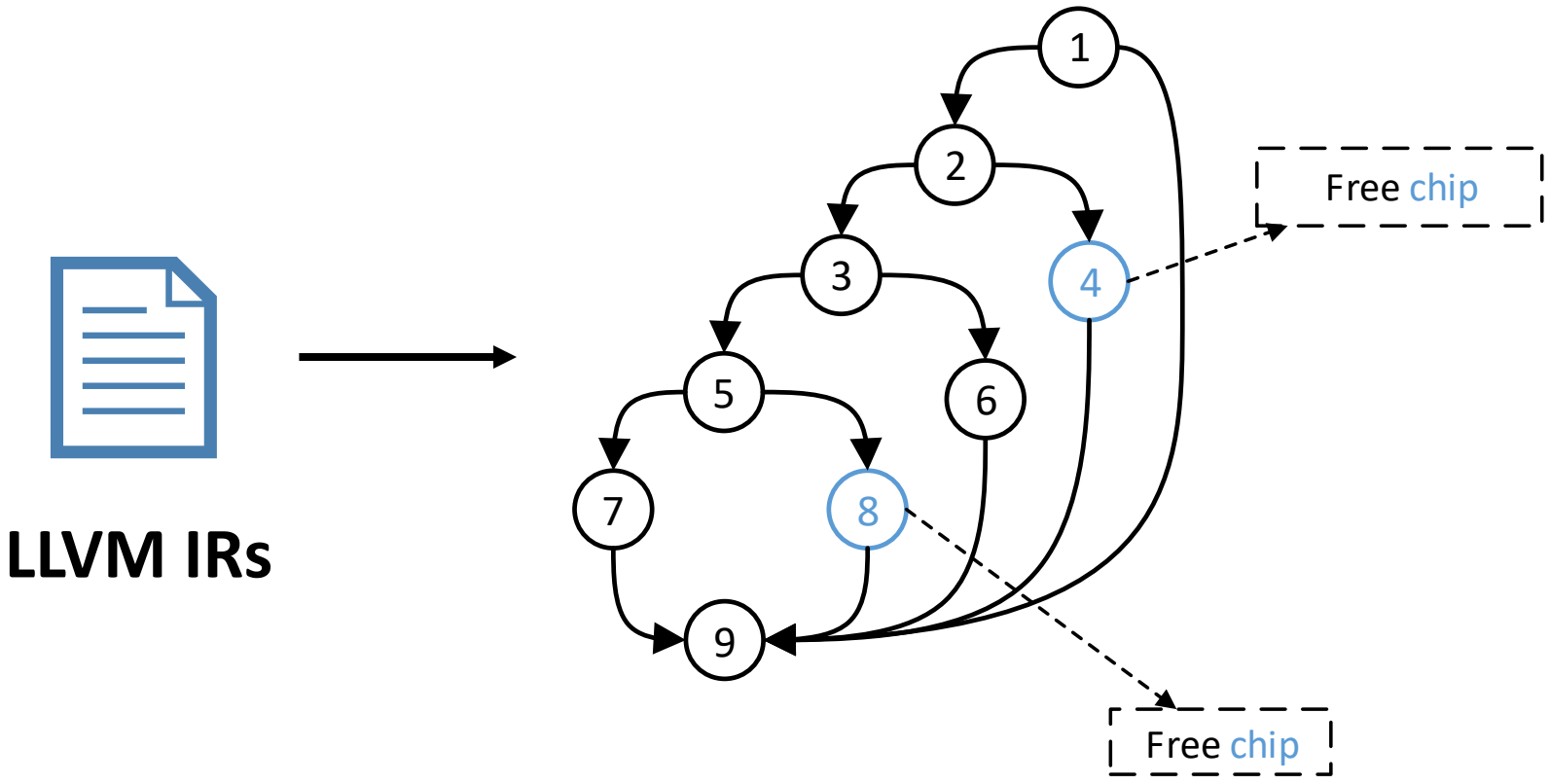
LLVM IRs



Control flow graph

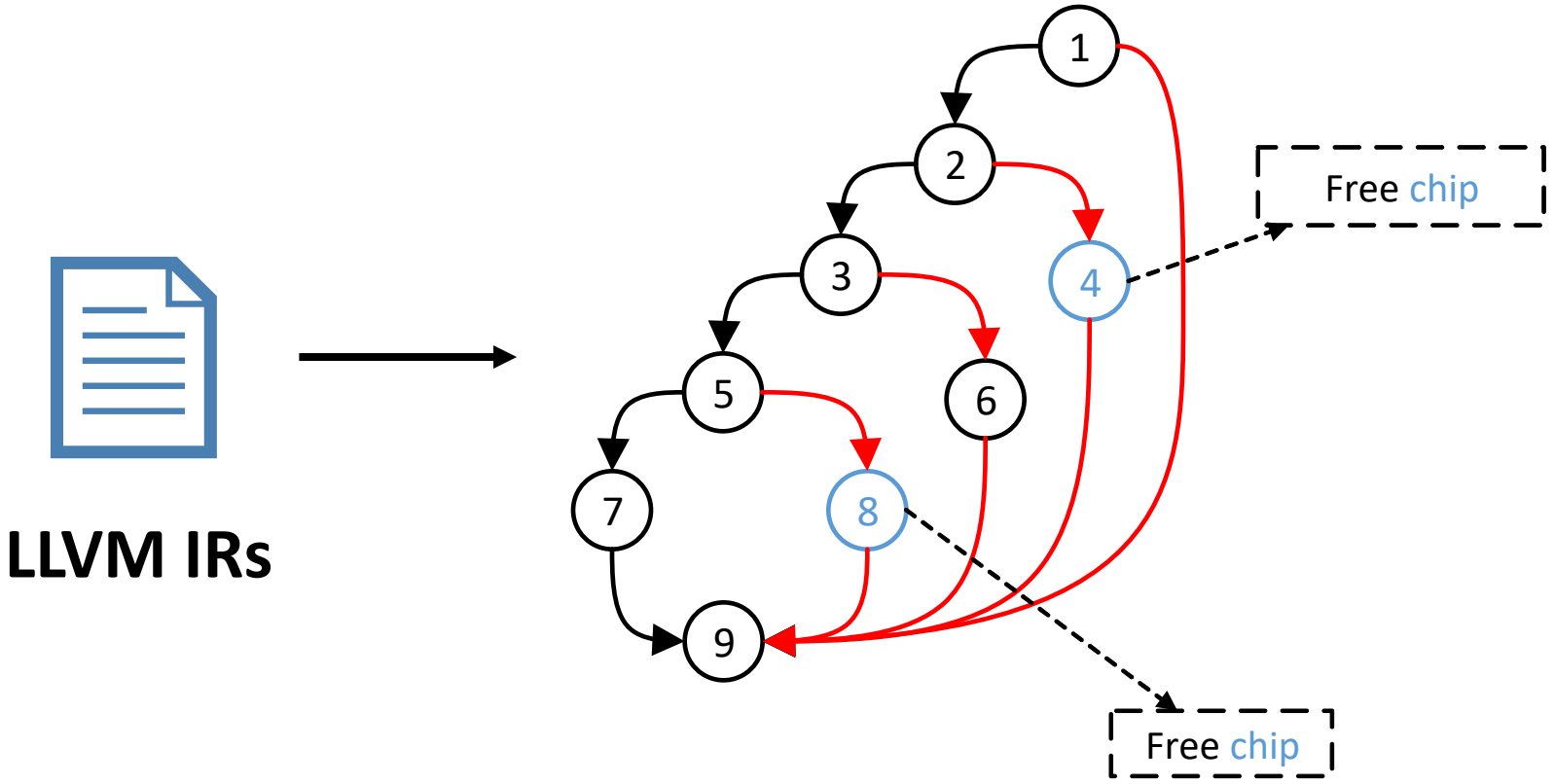
Workflow of IPPO

- **Security operation detection & error edges identification**



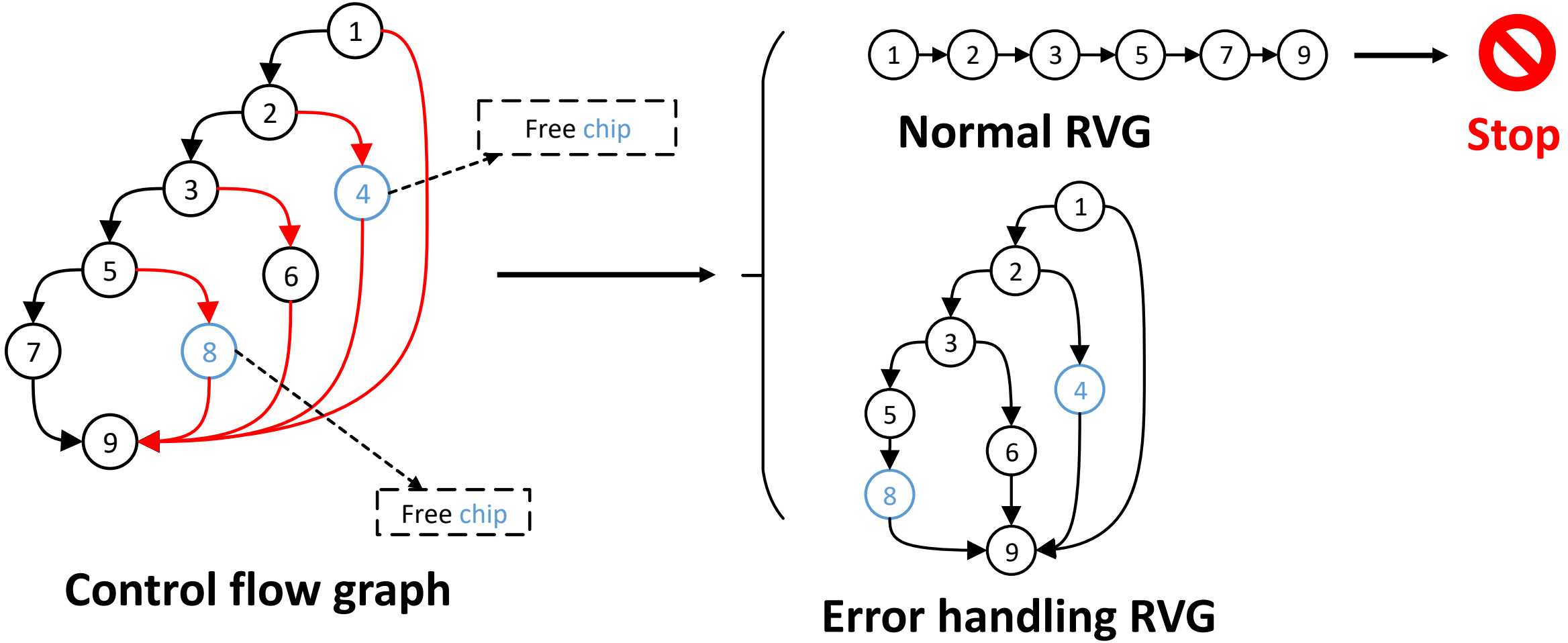
Workflow of IPPO

- **Identify error edges**



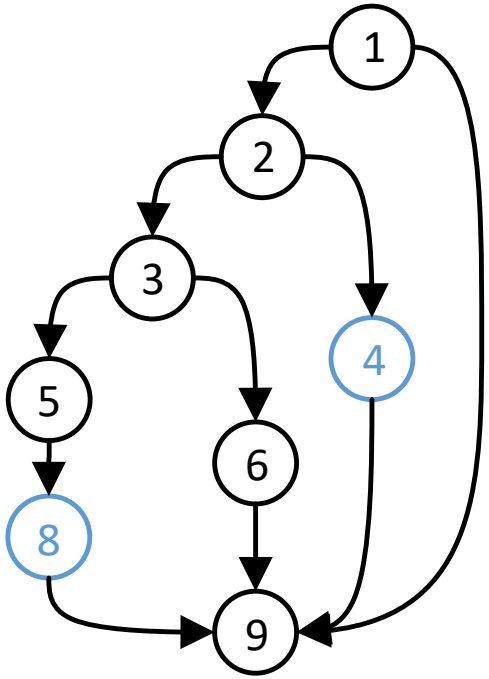
Workflow of IPPO

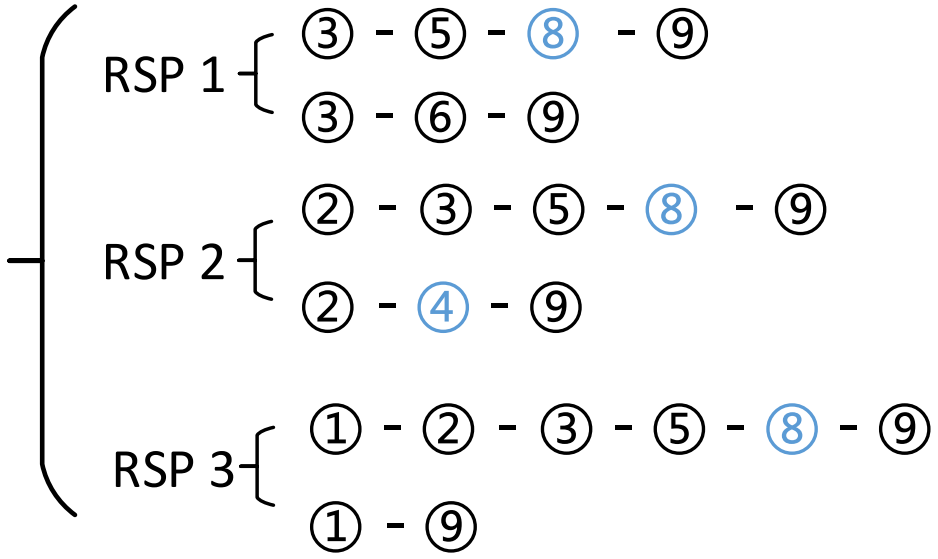
- **Generate return value-based graphs**



Workflow of IPPO

- **Collect reduced similar paths (RSPs)**



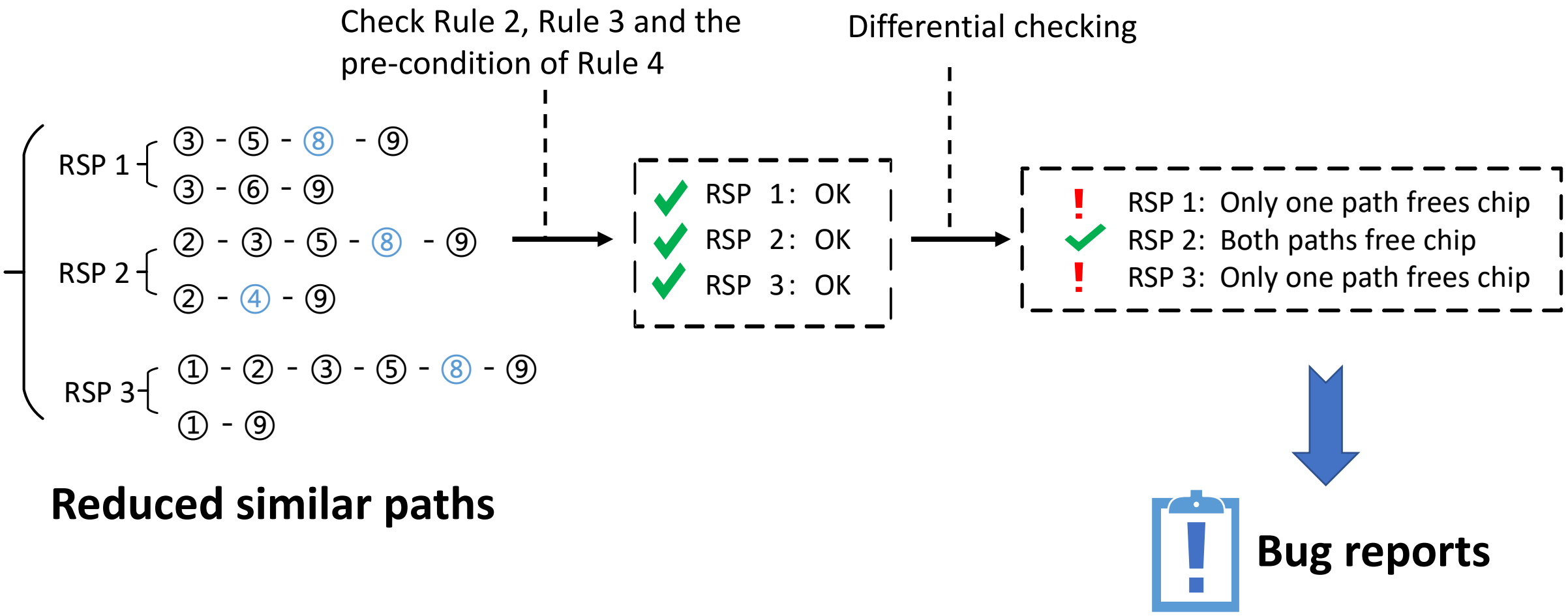


Error handling RVG

Reduced similar paths

Workflow of IPPO

➢ OSPP rules checking & differential checking



Evaluation

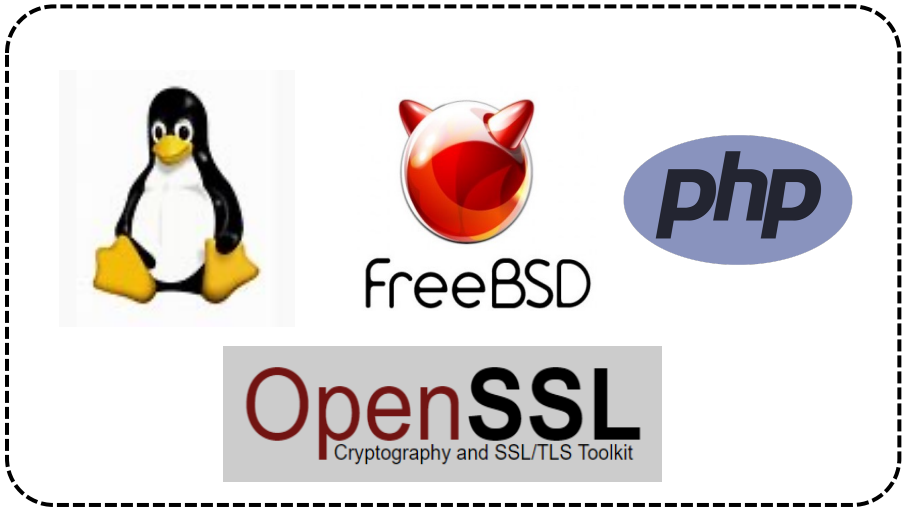
Experimental Setting

Environment

- Use a laptop with 16 GB RAM and Intel Core i7 CPU with six cores
- Use Clang-9.0

Targets

- Linux kernel v5.8
- FreeBSD 12
- OpenSSL 3.0.0-alpha6
- PHP 8.0.8



Bug Findings

- **Only focus on missed *return value checks*, *refcount decrement*, *resource release*, and *unlock*.**
- **Complete the whole analysis in 2 hours.**

Table 2: Bug detection results of IPP0 in the four systems. The R and T in the table indicate the reported bugs and true bugs, respectively.

Bug type	Linux		OpenSSL		FreeBSD		PHP	
	R	T	R	T	R	T	R	T
Missing check	101	11	2	1	1	0	4	0
Missing release	244	68	13	6	1	0	11	1
Refcount leak	345	181	0	0	0	0	0	0
Missing unlock	29	6	0	0	2	1	2	0
Total	719	266	15	7	3	1	17	1

275 valid bugs.

161 are previous unknown.

136 have been fixed by our patches or reports.

Comparison with Other Tools

➢ Comparison with cross-checking tools

Bug type	IPPO	FICS	Crix	APISan
Missing check	12	0	1	0
Missing release	75	0	0	0
Refcount leak	181	0	0	0
Missing unlock	7	0	0	0
Total	275	0	1	0

IPPO is a promising complementation with existing tools.

➢ Comparison with pairing analysis tool: HERO

Bug types	**Bugs in v5.3**	**HERO Results**	**Recall**
Memory Leak	55	2	3.6%
Refcount Leak	112	82	73.2%
Missing unlock	3	0	0%
UAF/DF	6	0	0%
Total	176	84	47.7%

Limitation & Discussion

➢False positives

- Unexpected pre-condition.
- Imprecise data-flow analysis.
- Imperfect error path analysis.
- Imperfect security operation detection.
- ……

➢False negatives

- Imperfect security operation detection.
- ……

➢Supporting inter-procedural analysis

- Model inter-procedural object-based similar paths.

Conclusion

- **Missing security operations is common in real-world programs, and could cause various security issues.**

- **We presented IPPO: a framework to detect missed security operations.**
  - Object-based similar path pairs.
  - Reduced similar path.
  - Return value-based sub-CFG.

- **We evaluated IPPO on 4 real-world programs.**
  - Find 161 new bugs.
  - IPPO could effectively detect bugs that missed by existing tools

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